

A Brief Analysis of the SPEC CPU2000 Benchmarks on the Intel® Itanium® 2 Processor

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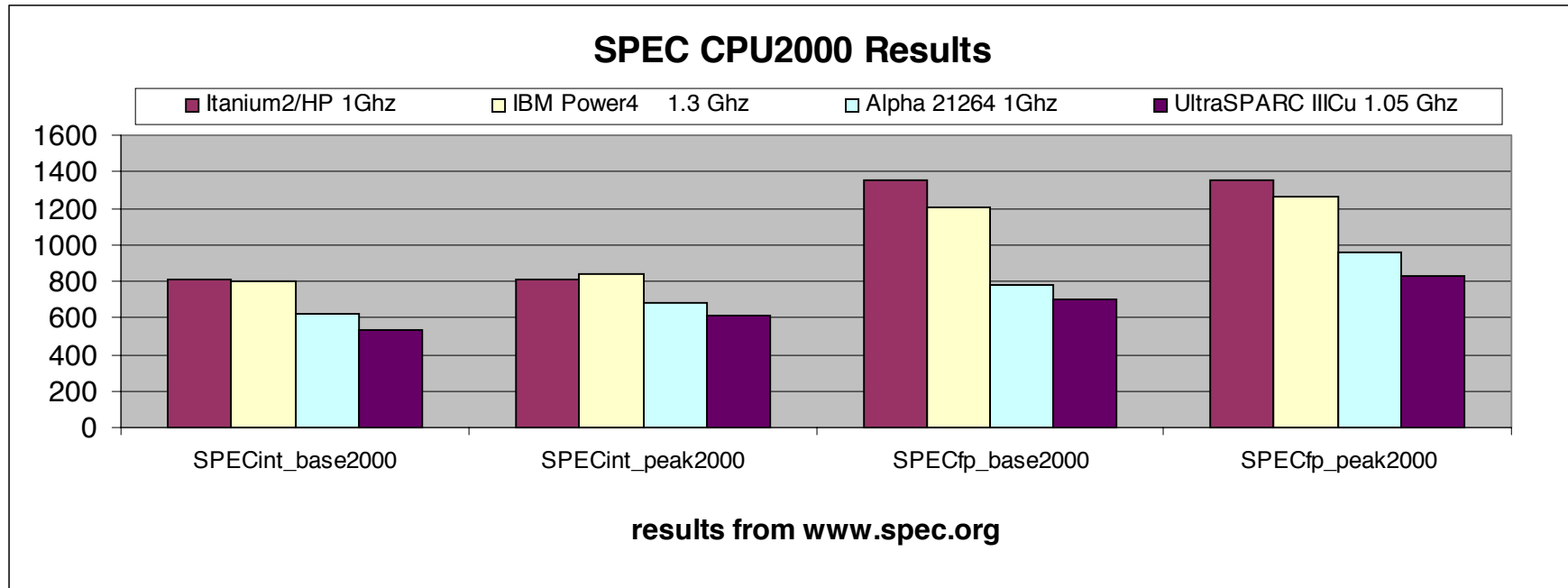
Agenda

This is a data-oriented presentation, not research

Agenda

- Brief performance summary
- Comparison of how HP and Intel compilers use Itanium® architecture features and compare to best RISC.
- Analysis of microarchitectural features of the Intel® Itanium® 2 processor and how it affects SPEC CPU2000 performance

Overall Performance



SPEC{int/fp}_base2000

810/1356 (best 0.18u)

Linpack 1000:

3.5 Gflops (best overall)

TPC-C (SQL/4P):

78K tpmC (best 4P number)

SPECweb_SSL:

1520 connections (best of class)

➤ **Itanium® 2 processor best of class on a wide range of applications**

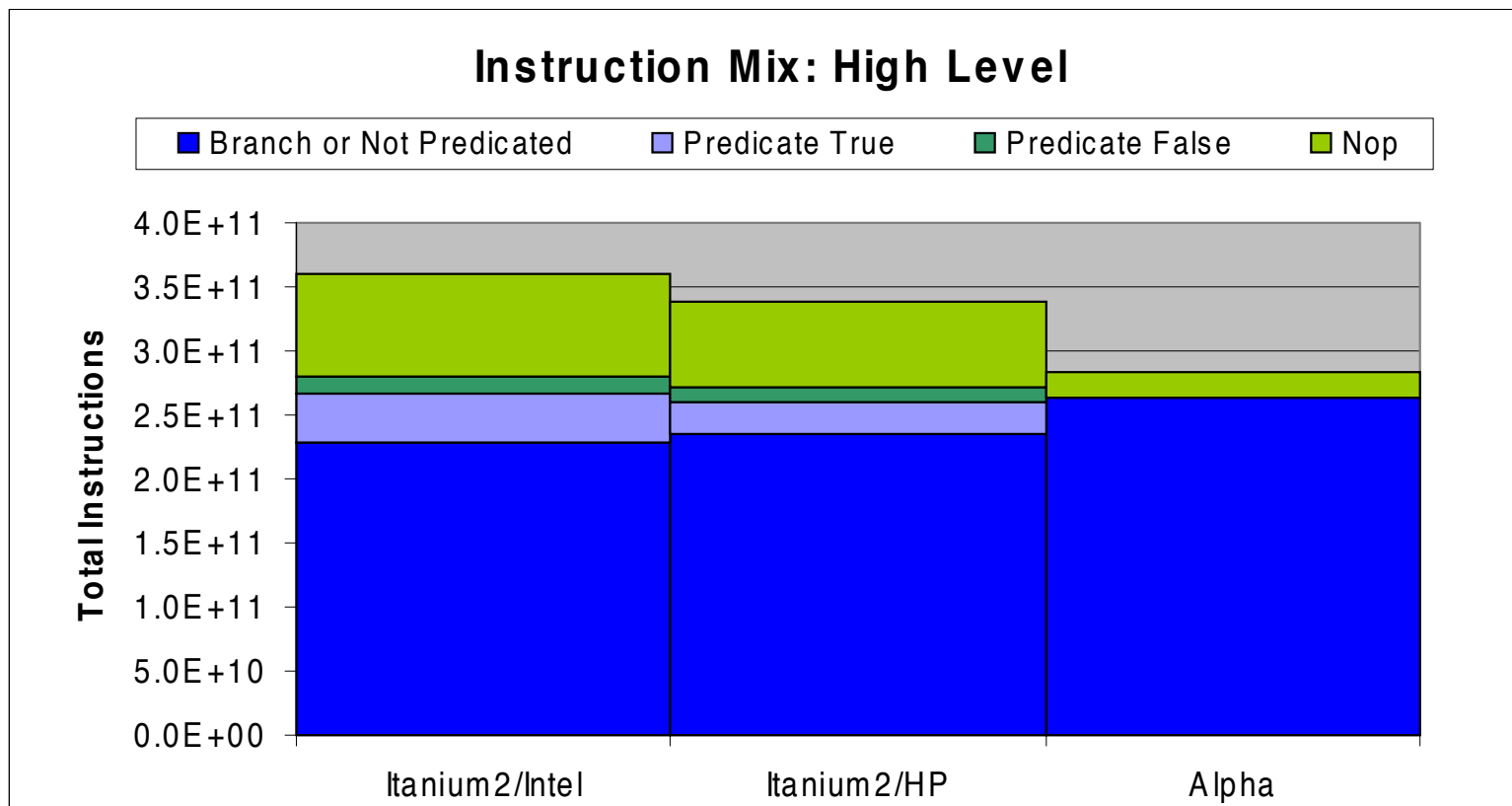
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Part I: ISA and Compiler Comparisons

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Team

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- Number of ‘useful instructions’ (shown in blue) +/- 1% of Alpha
 - Total instructions (blue+green) +20-30% of Alpha (due to NOPs)
 - Bundling is main cause of extra NOPs
- **We’ll see that extra instr/nops are not substantially impacting perf**

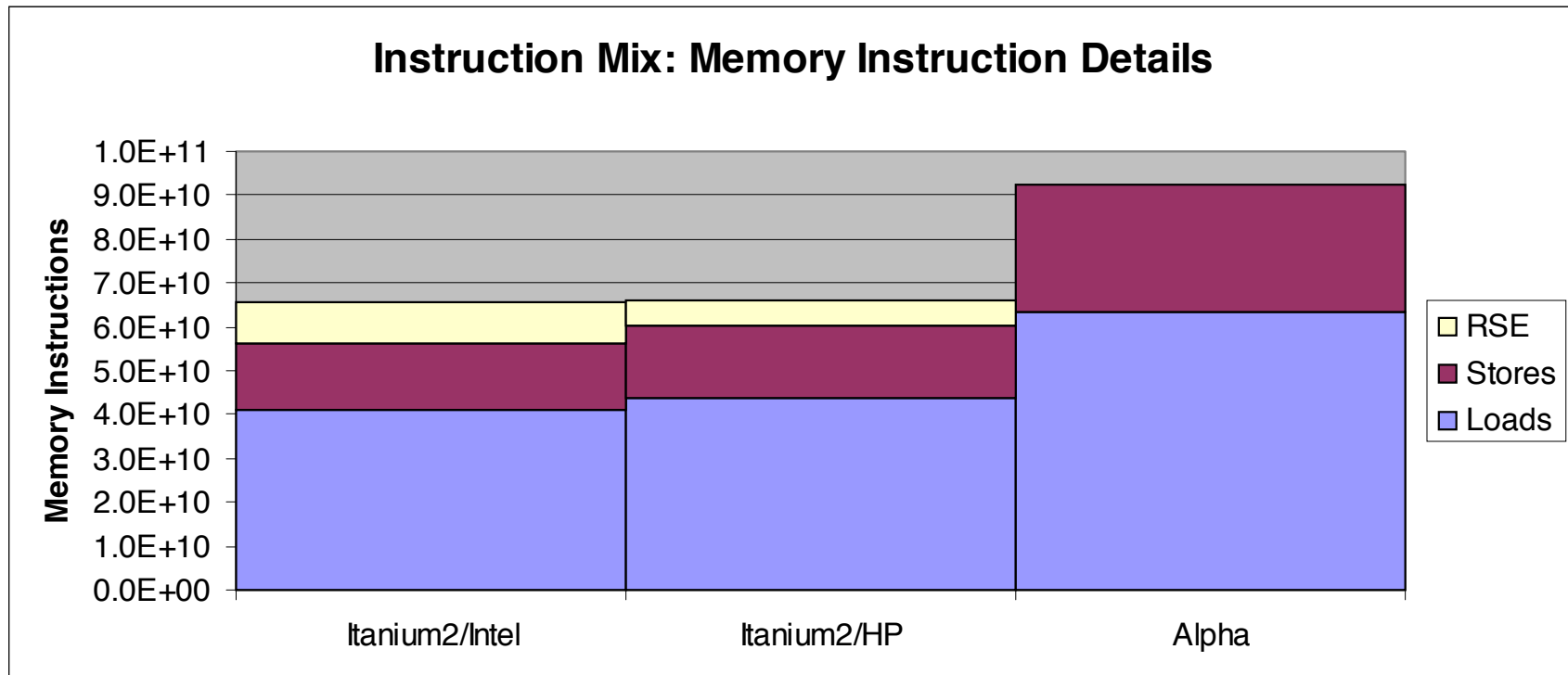
Instruction Mix Details Introduction

Itanium/Alpha ISAs:

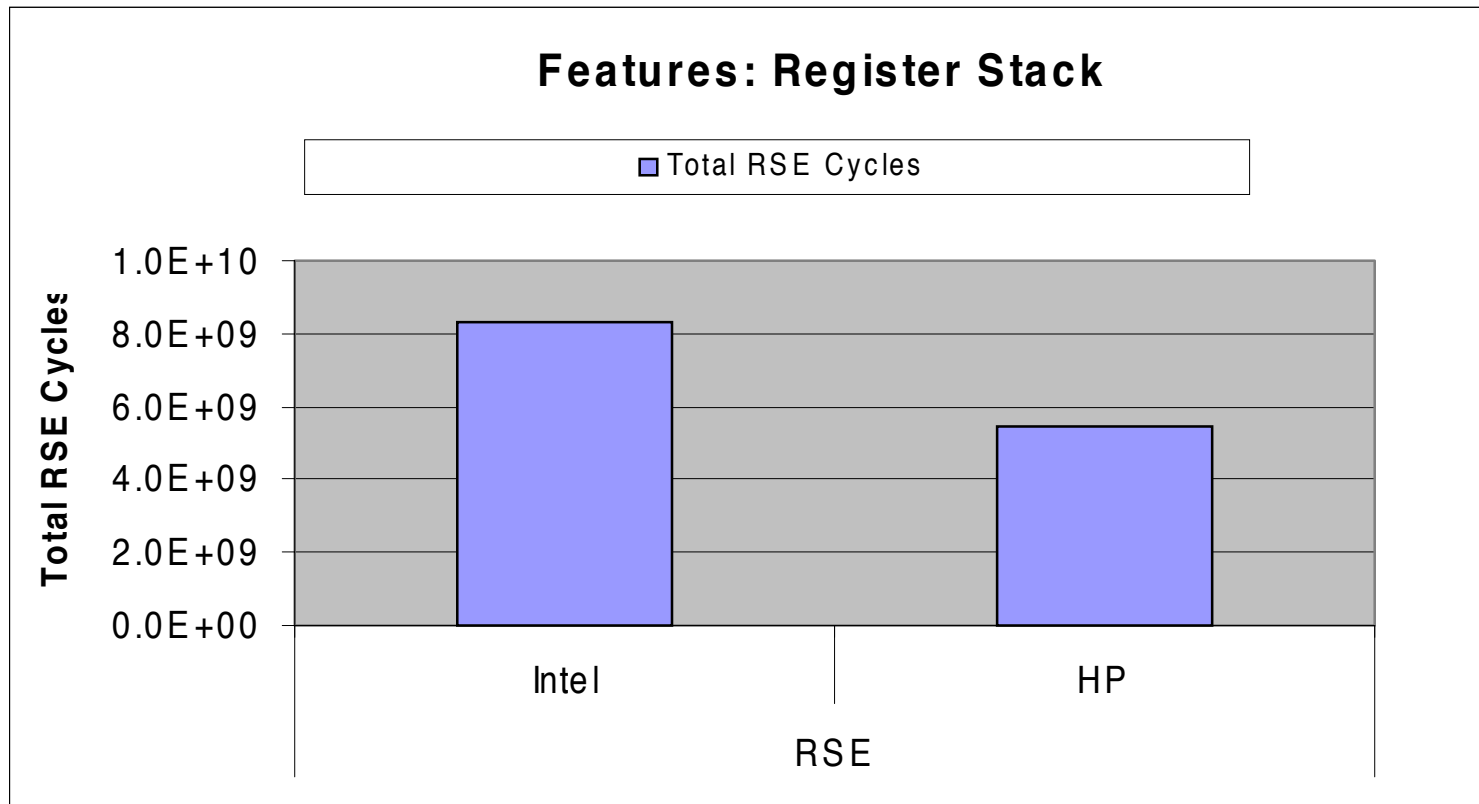
- Itanium® arch has 40% fewer memory operations and 30% fewer branches than Alpha (some impact from no-pgo on Alpha)
- Itanium arch has about 10% more ALU/compares/shifts than Alpha
- Itanium arch has about 20-30% NOPs – eventually expect this to be under 20%

HP/Intel Compiler:

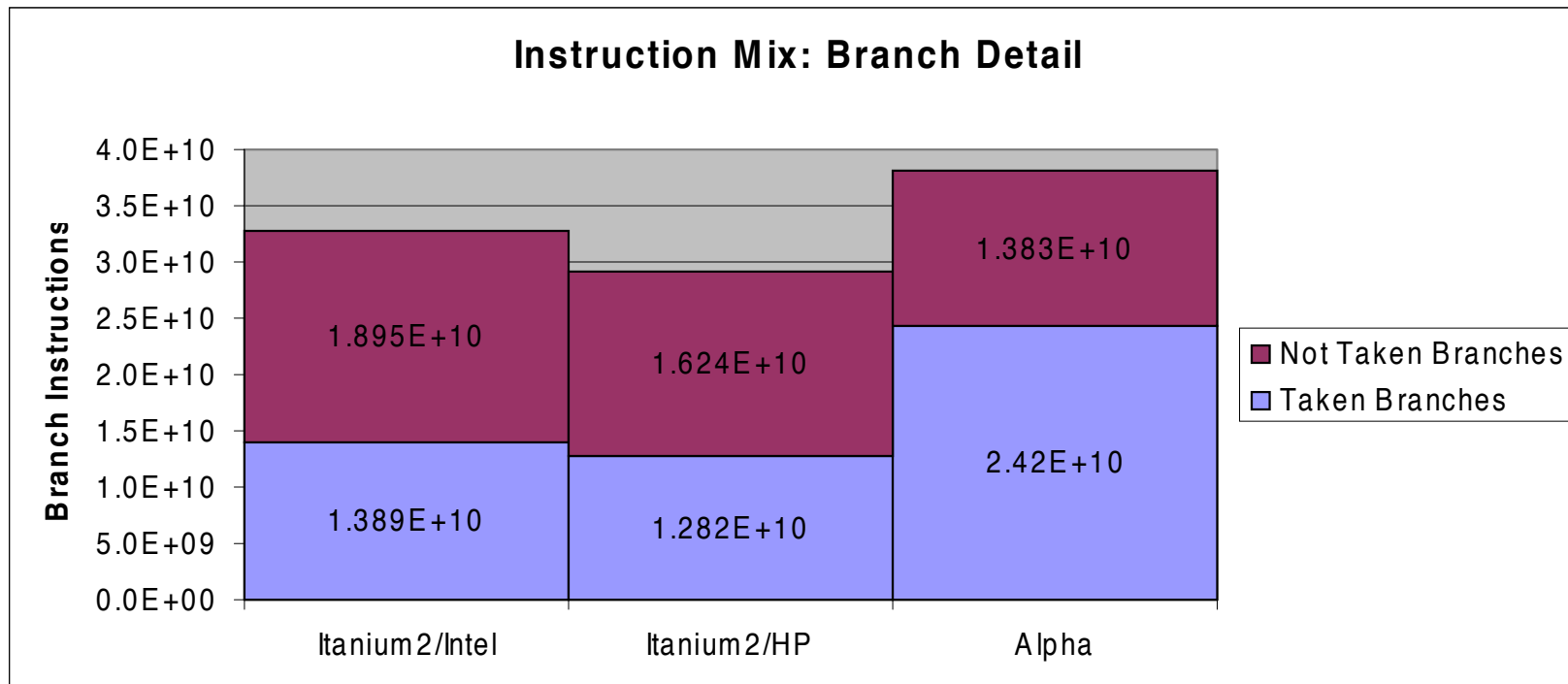
- HP compiler uses more memory and ALU ops than Intel
- Implies HP more conservative with registers – we'll see impact later
- **Itanium® architecture trades more 'easy' instructions (NOP, alu, cmp) for reducing the 'hard' instructions (branch, load)**
- **More than one way to get good performance from a compiler**



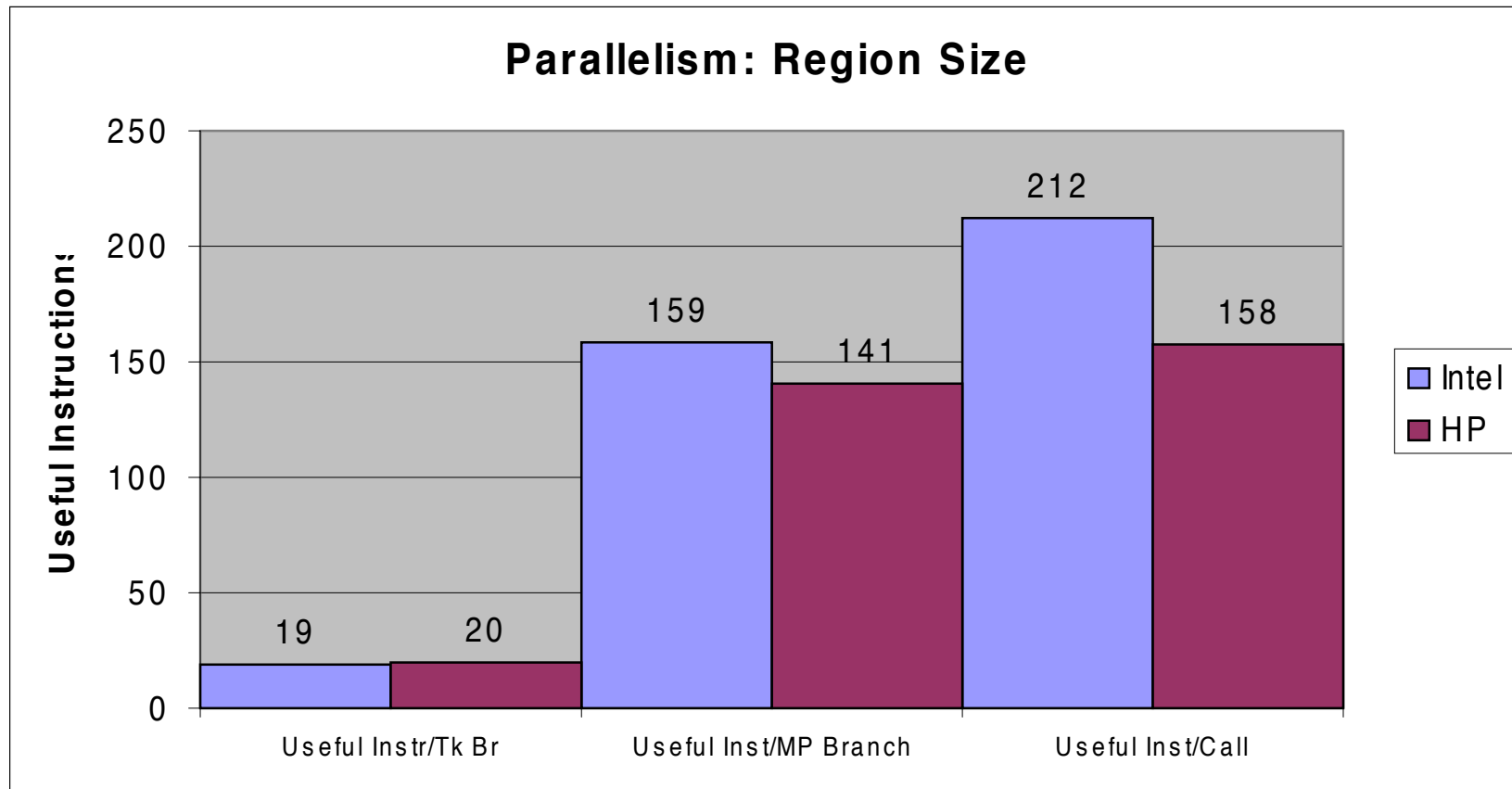
- Alpha has 1.4x mem refs of Itanium® architecture (incl/RSE ops)
 - RSE ops are easy to optimize in the future, if needed
 - HP compiler is more consv with regs, but has more ALU/reloads
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- **Large register file, good compiler technology, and RSE pay off**
 - **HP has lower RSE costs, but more reloads/alu ops**



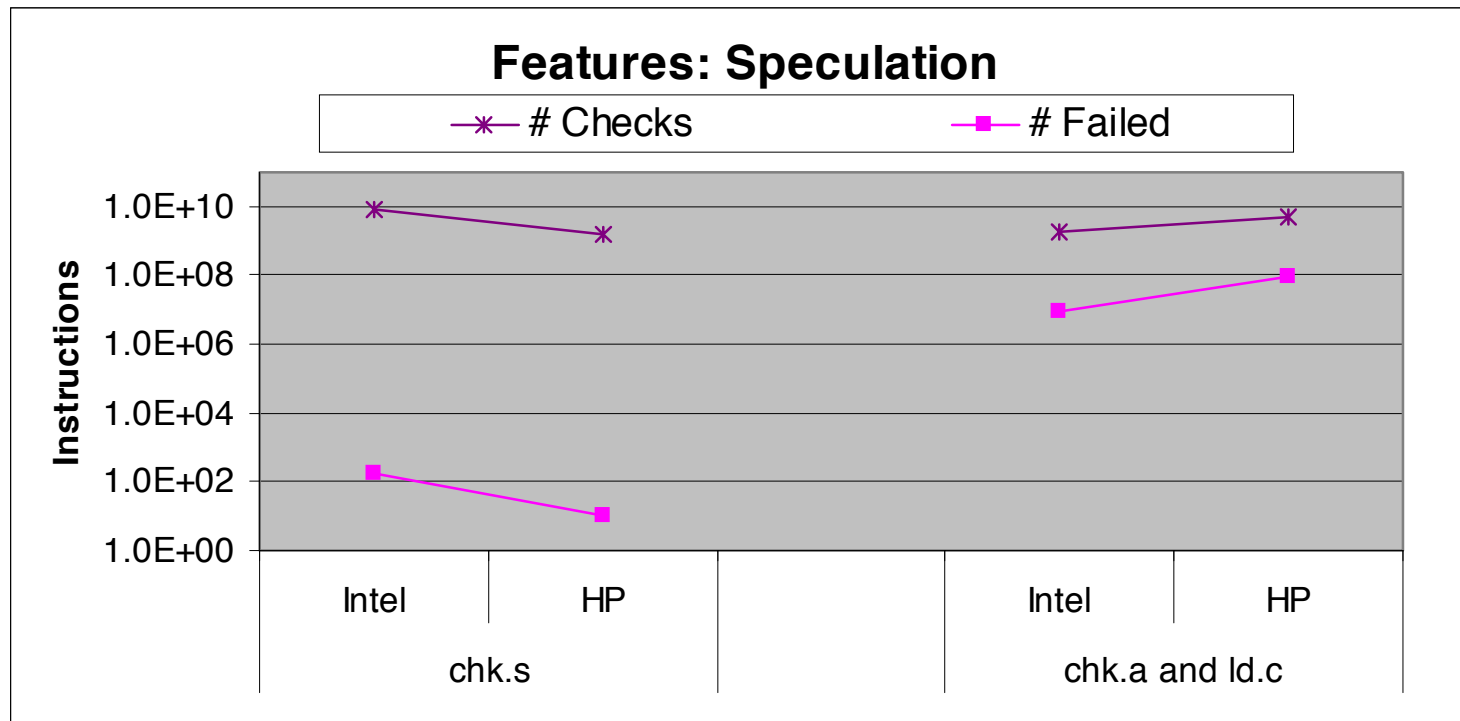
- Total time spent in RSE is only 3%-4% of overall execution
 - 1½ -3 cycles per call/return for RSE spill/fill activity
 - 1-3 instructions per subroutine setup for RSE
 - Intel compiler has fewer calls, but more cycles/call
- **Register stack provides very low overhead call/return support**



- HP compiler generates 13% fewer br's than Intel, 9% more mispredictions
 - HP compiler generates 31% fewer branches than Alpha
 - Itanium-based binaries fewer tk br's than Alpha, data skewed by lack of PGO
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- **Itanium® architecture reduces the # branches and branch mispredicts**
 - **HP/Intel compilers both reduce # branches – but with different focus**

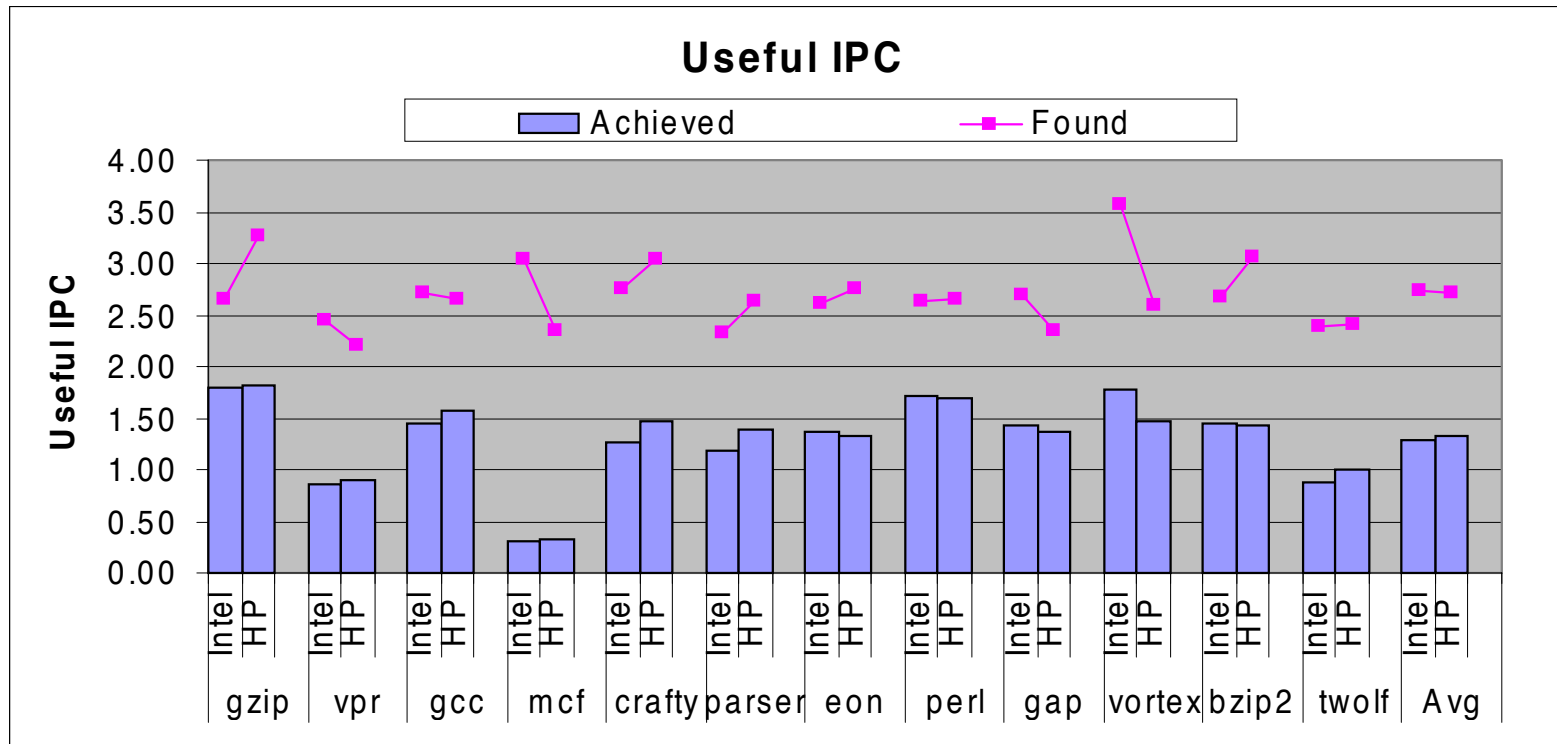


- Useful instructions per taken branch is very high
 - Useful instructions per mispredicted branch slightly better for Intel
 - Useful instructions/call very high – Intel/HP compilers very aggr inlining
- **Advanced compilers reduce stress on br prediction/Icache HW**
- **Trading Istream size for regularity improves HW efficiency**



- About 20-30% of loads are speculative in Intel binaries
- Data shows tiny penalty for chk.s usage despite high usage rate
- Intel has 10x more chk.s than HP, HP uses 'no recovery model' selectively (per benchmark decision)
- HP has 10x more chk.a/ld.c then Intel, recovery less than 1% time

➤ **Speculation heavily used, but causes little overhead**



- Useful IPC computed using ‘unstalled IPC’
 - Compilers find 2.5-3.0 IPC in integer apps (even beyond SPEC)
 - Dynamic delays reduce this to 1.3 achieved for CPU2000 integer
-
- **Differences in perf/heuristics shows headroom for both compilers**
 - **Good IPC found by both compilers, room for uArch improvements**

Notes

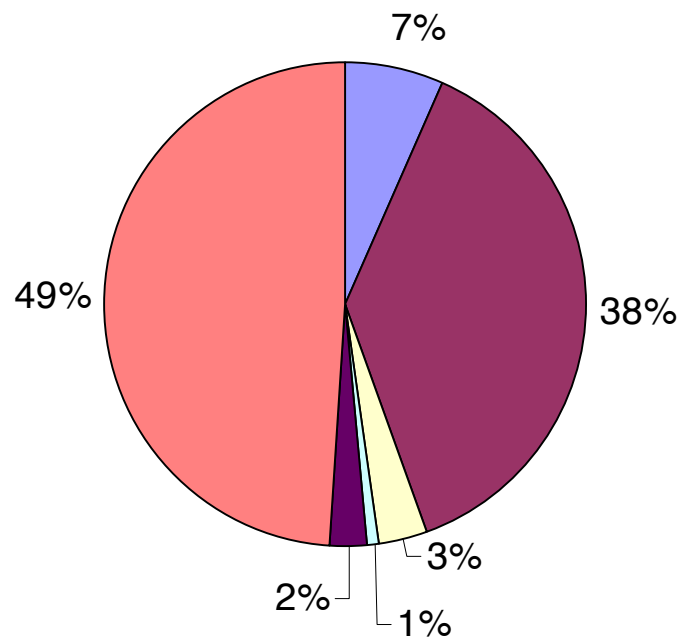
- Itanium-based binaries used for these stats are older than those used for the official SPEC submission (less than 10% difference)
- The results for Intel® Itanium® 2 processor in Part I are: one with the Intel compiler running on 64-bit MS OS and another with the HP compiler running under HP-UX
- Alpha ISA numbers via simulation, binaries used near peak (no profile guided optimization), tuned for 21264. Alpha data missing VPR and PERL – thus, left out of all averages in Part I.
- Results computed with arithmetic averages – data thus skewed towards long-running benchmarks

Part II: Microarchitecture

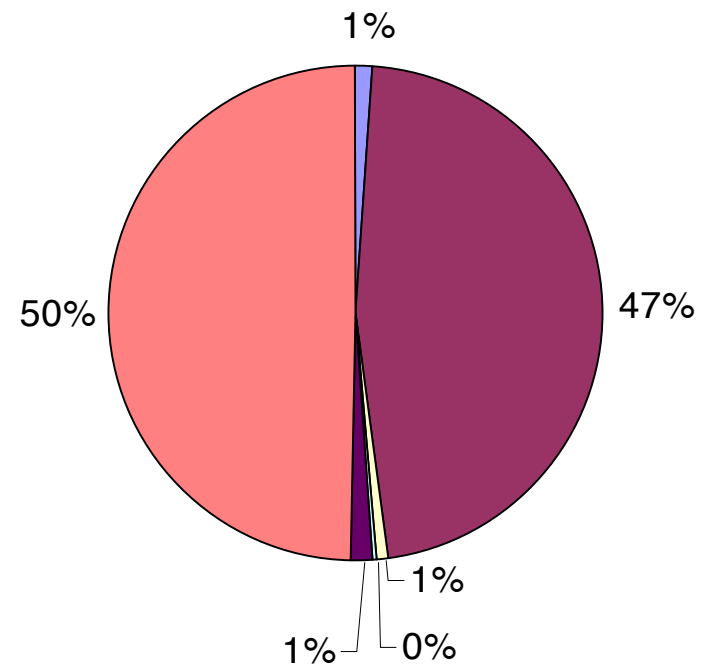
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Cycle Accounting: INT

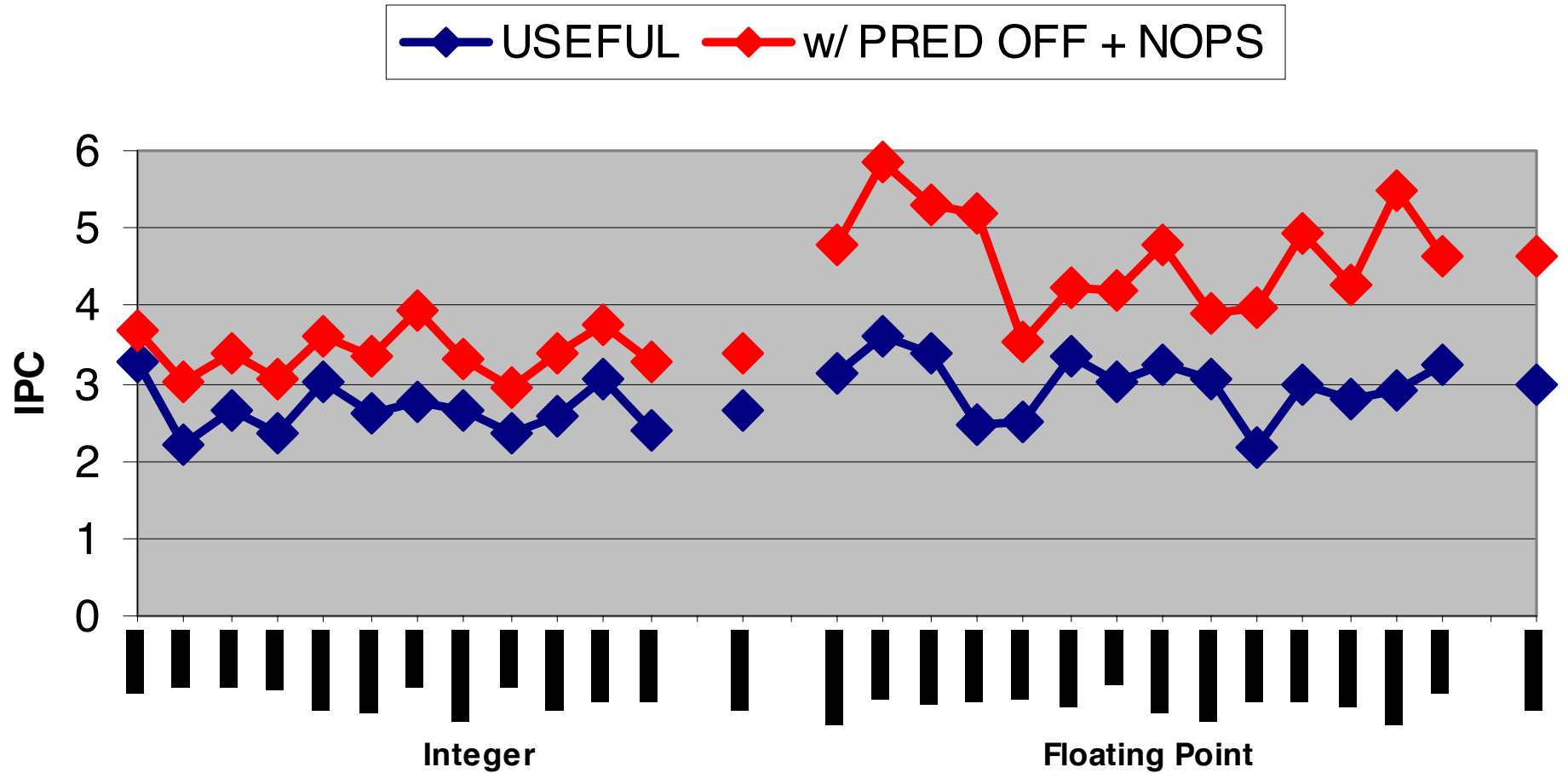


Cycle Accounting: FP



- Remaining: *unstalled execution* and *data access*
- Great performance

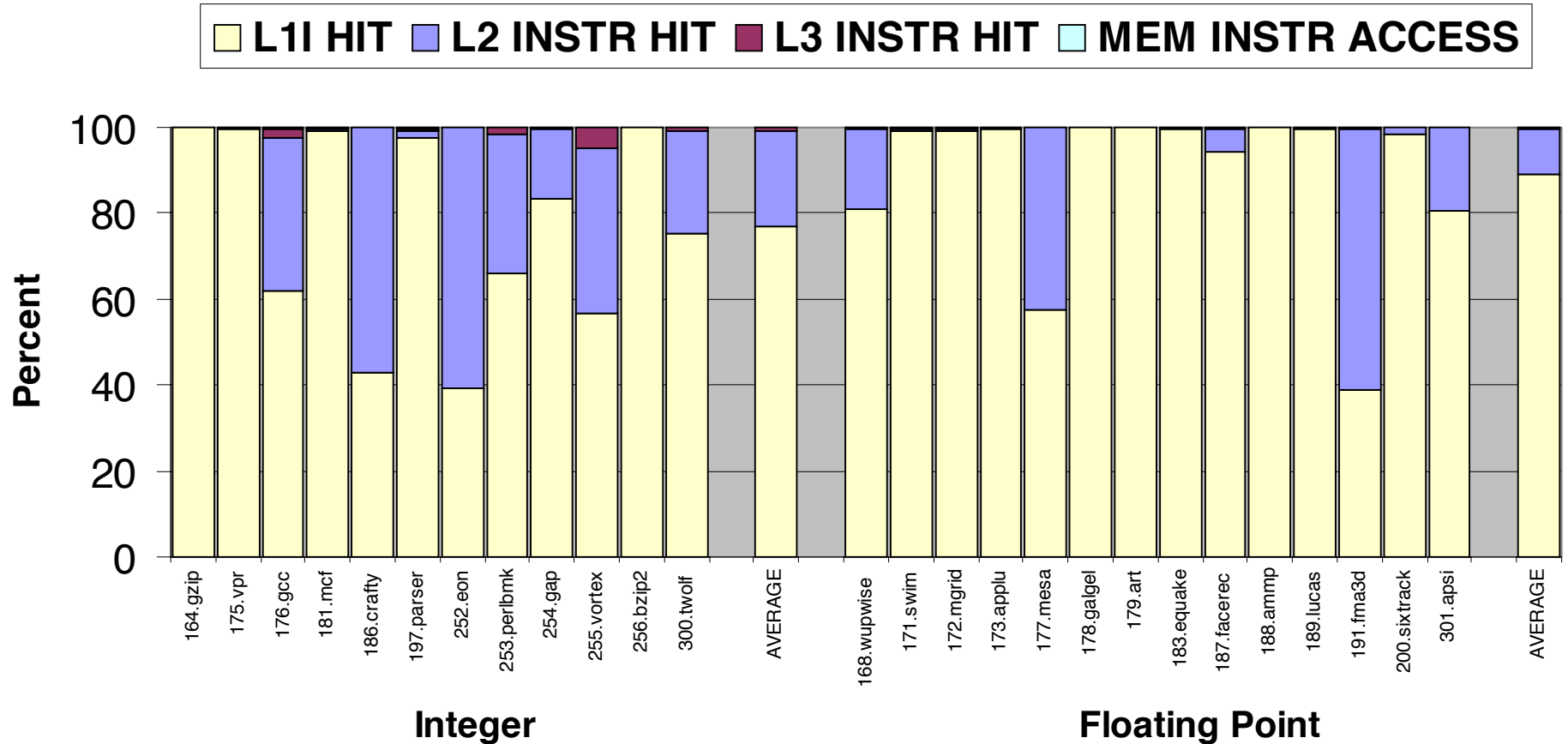
Instructions Per Unstalled Cycle



➤ High machine parallelism

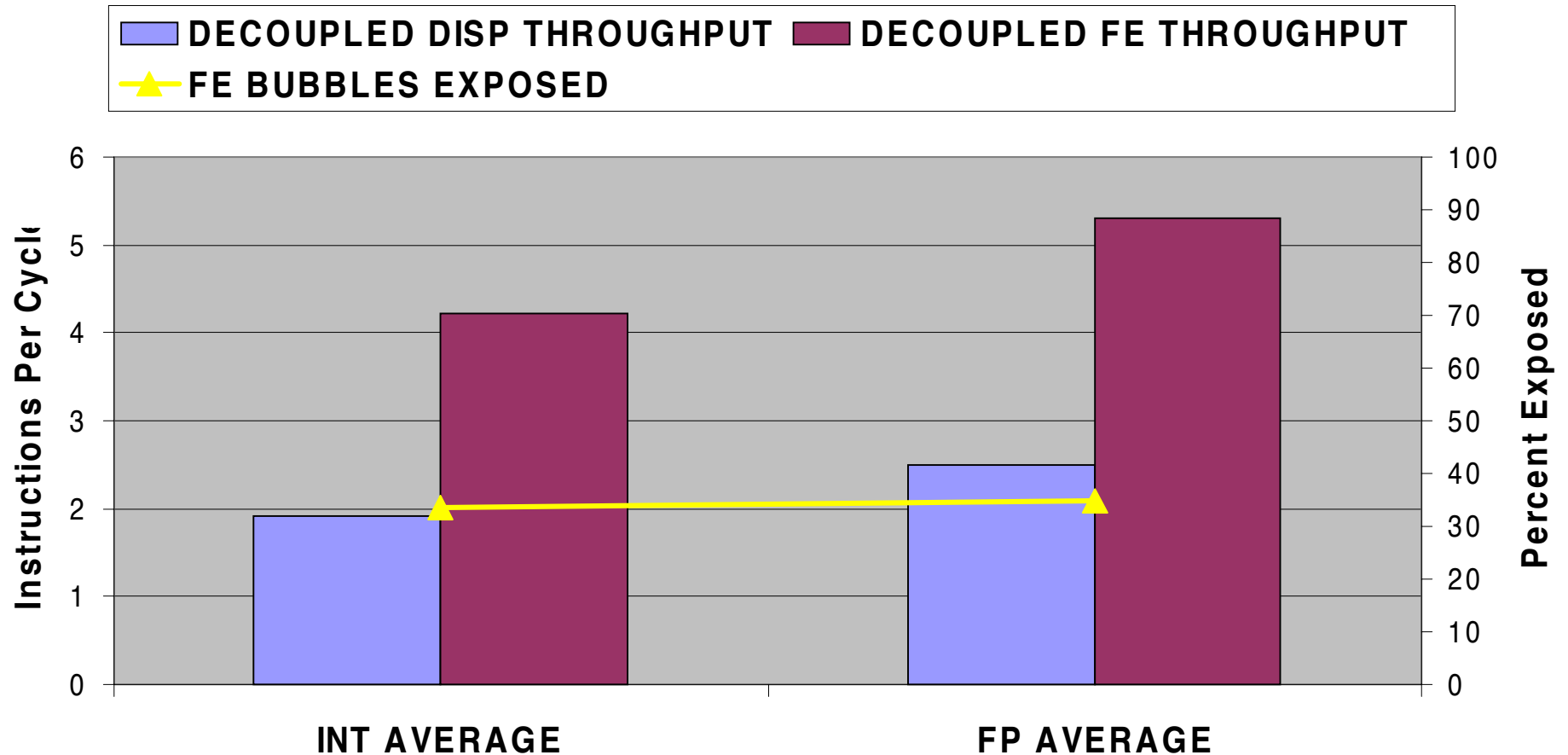
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Total Instruction Fetch Latency



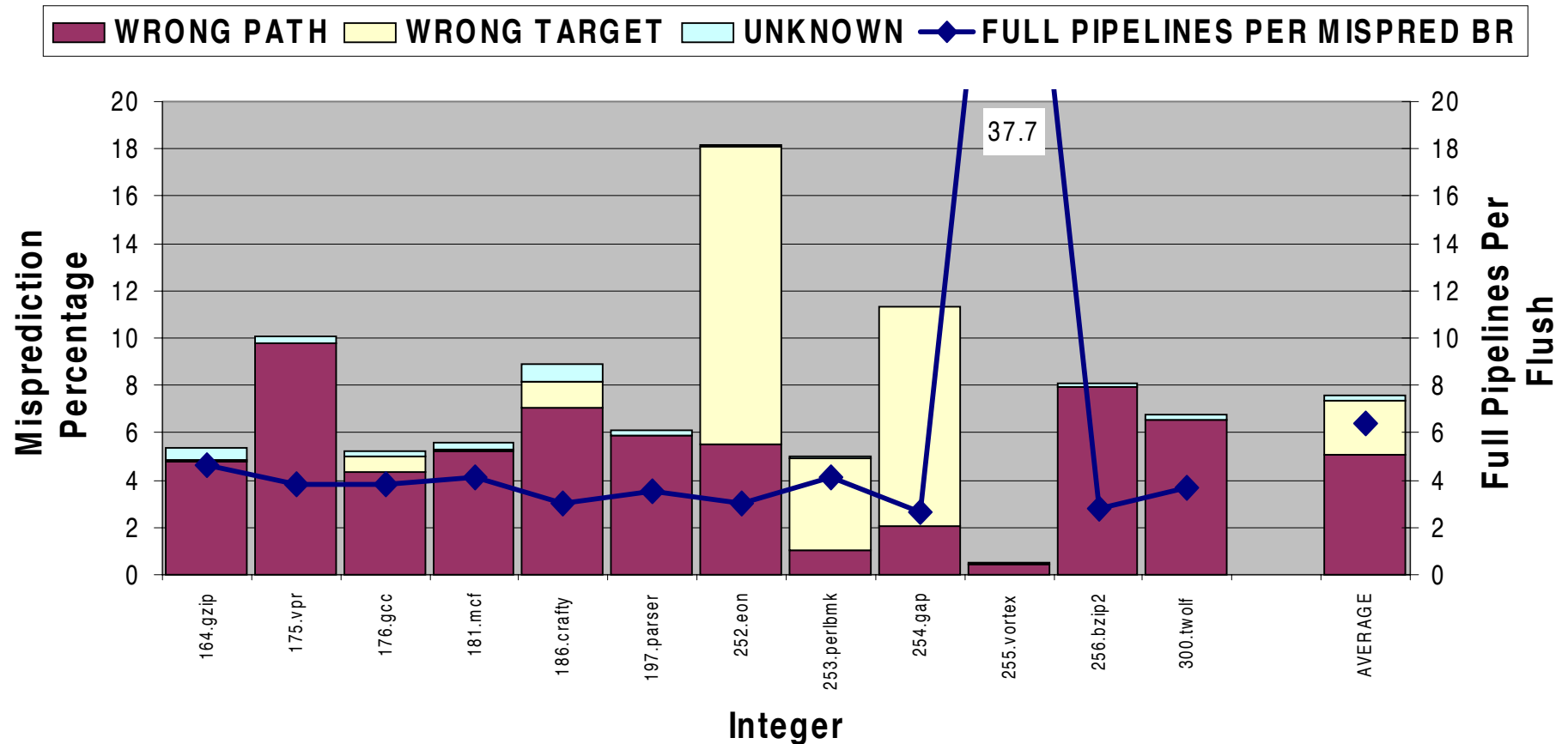
- Noticeable L1I misses
- Very small I-fetch component

Instruction Fetch Ahead



- High FE throughput rate
- I-miss latency hidden

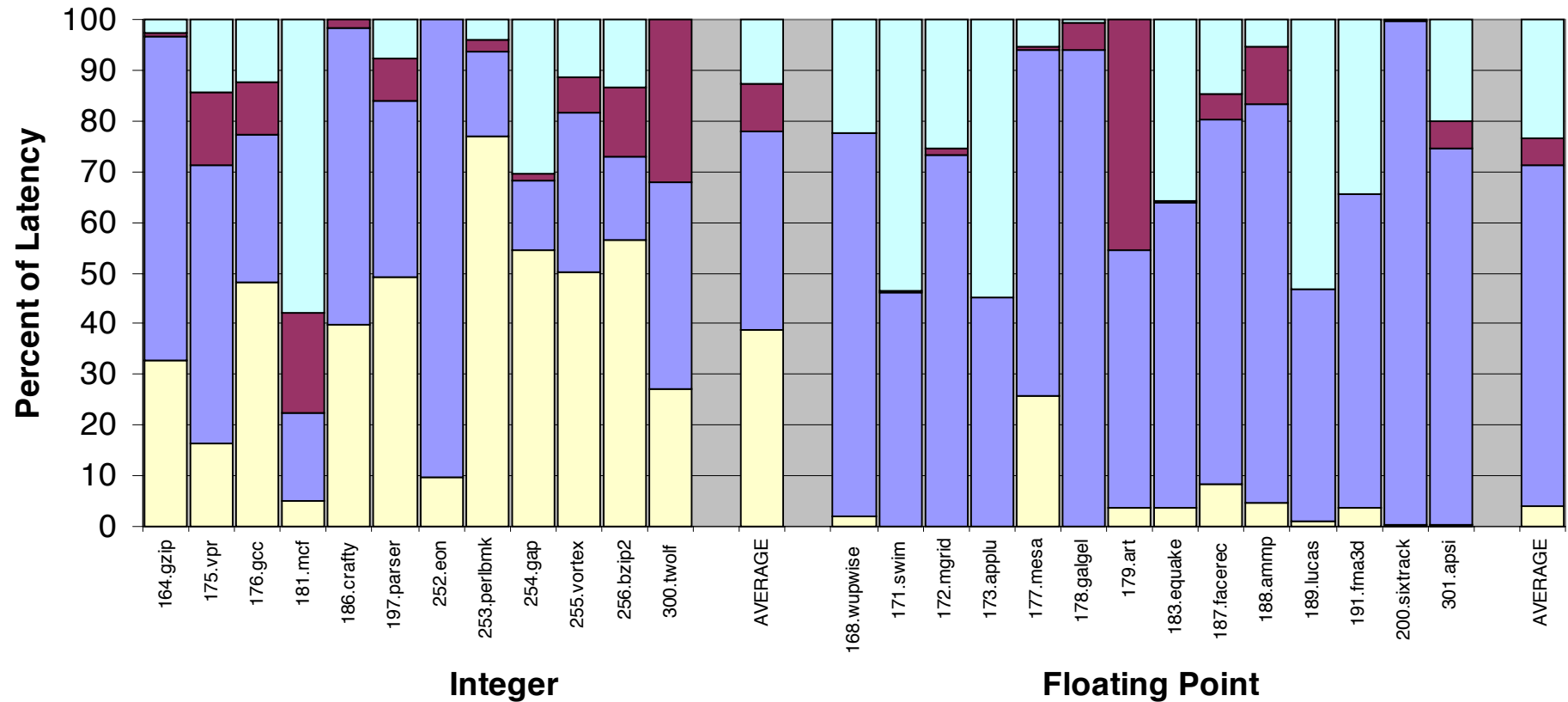
Branch Mispredictions



- High accuracy, low penalty
- Helps instruction fetch

Total Data Read Latency

■ L1D DATA READ
 ■ L2 DATA READ
 ■ L3 DATA READ
 ■ MEM DATA READ



➤ Large component, large opportunity

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Summary

Itanium® 2 Processor Delivers Leadership Performance

- Architecture / Compilers
 - Expose lots of ILP to in-order pipeline
 - Replace difficult instructions with easy ones
 - RSE and large register file work well together
- CPU Design
 - Machine parallelism handles high ILP
 - Aggressive design reduces bottlenecks

Acknowledgements

Jason Cantin (U. Wisconsin):

Ran all the experiments to gather the Alpha ISA data with changes/alterations at our request. Without Jason, we would have no Alpha data of any kind. Thanks to his efforts to rerun all the data for us!

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Bryan Black, Ed Grochowski, Jim Callister, Carole Dulong (Intel):

Extensive draft review and comments

Caliper Team (HP):

Tool support

Backup/Reference Slides

Configuration Information

Intel® Itanium® 2 processor data for Intel systems/compiler:

- Binaries: Pre-production version of Intel C++ 6.0 Compiler, -O3 with interprocedural optimization and profile guided optimization
- Run on: Pre-production stepping of Itanium 2 processor 800Mhz/200Mhz core/bus, Intel 870 chipset, monitor kernel and user instructions and events

Intel® Itanium® 2 processor data for HP systems/compiler:

- Binaries: Pre-production version of HP compiler
- Run on: Pre-production stepping of Itanium 2 processor 1000Mhz/200Mhz core/bus, rx2600 prototype, monitor kernel and user instructions and events

Alpha ISA data:

- Run on: Functional simulator, system code not simulated
- Simulation details: <http://www.cs.wisc.edu/multifacet/misc/spec2000cache-data/>
- All data thanks to Jason Cantin at the University of Wisconsin.
- Binaries: <http://www.eecs.umich.edu/~chriswea/benchmarks/Com.cf>
In general, these binaries are optimized for 21264, peak optimization. Usually, -g3 -fast -O4, but NO profile feedback. Compiler: DEC C V5.9-005 and DIGITAL C++ V6.1-027

Other remarks:

- All averages in slides left out PERL and VPR due to data not available for Alpha