

# picoJava<sup>™</sup>: A Hardware Implementation of the Java Virtual Machine

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# The Java – picoJava Synergy

- Java's origins lie in improving the consumer embedded market
- picoJava is a low cost microprocessor dedicated to executing Java<sup>™</sup>-based bytecodes

-Best system price/performance

### It is a processor core for:

- -Network computer
- -Internet chip for network appliances
- -Cellular phone & telco processors
- -Traditional embedded applications



# **Java in Embedded Devices**

### Products in the embedded market require:

### Robust programs

-Graceful recovery vs. crash

### Increasingly complex programs with multiple programmers

-Object-oriented language and development environment

### Re-using code from one product generation to the next

-Portable code

### Safe connectivity to applets



-For networked devices (PDA, pagers, cell phones)

# Important Factors to Consider in the Embedded World

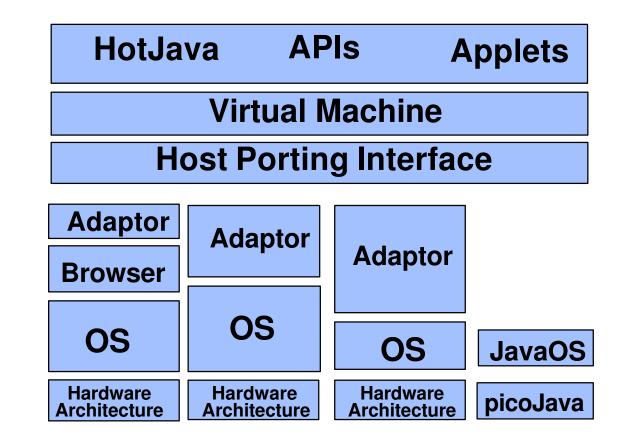
### Low system cost

-Processor, ROM, DRAM, etc.

- Good performance
- Time-to-market
- Low power consumption



# Various Ways of Implementing the Java Virtual Machine





# picoJava

### Directly executes bytecodes

- -Excellent performance
- Eliminates the need for an interpreter or a JIT compiler
- -Small memory footprint

### Simple core

-Legacy blocks and circuits are not present

### Hardware support for the runtime

-Addresses overall system performance



# **Java Virtual Machine**

#### • What the virtual machine specifies:

- -Instruction set
- -Data types
- -Operand stack
- -Constant pool
- -Method area
- -Heap for runtime data
- -Format of the class file



# Virtual Machine —Instruction Set

- Data types: byte, short, int, long float, double, char, object, returnAddress
- All opcodes have 8 bits, but are followed by a variable number of operands (0, 1, 2, 3, ...)
- Opcodes
  - -200 assigned
  - -25 quick variations
  - -3 reserved



# **Java Virtual Machine Code Size**

#### ■ Java<sup>™</sup>-based bytecodes are small

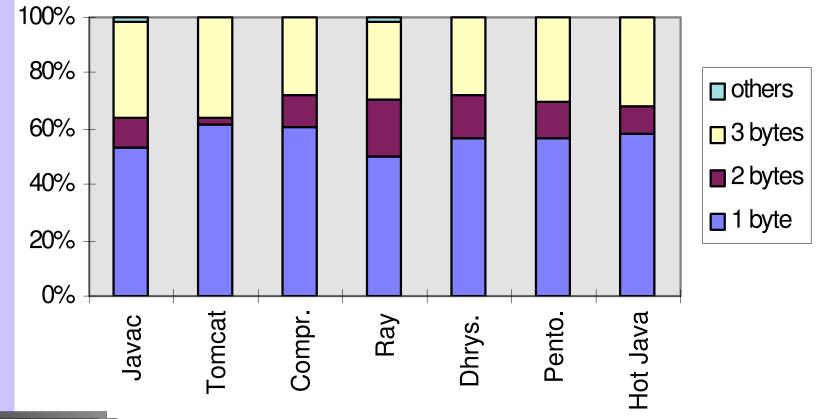
- -No register specifiers
- -Local variable accessed relative to a base pointer (VARS)

### This results in very compact code

- -Average JVM instruction is 1.8 bytes
- -RISC instructions typically require
  - 4 bytes



# **Instruction Length**





# **Java Virtual Machine Code Size**

- Java bytecodes are about 2X smaller than the RISC code from the C++ compiler
- A large application (2500+lines) coded in both the C++ and Java languages



# JVM – Instruction Set – RISCy

### Some instructions are simple

bipush value	:push signed integer
iadd	integer add:
fadd	single float add:
ifeq	:branch if equal to O
iload offset	:load integer from
	:local variable



# JVM – Instruction Set – CISCy

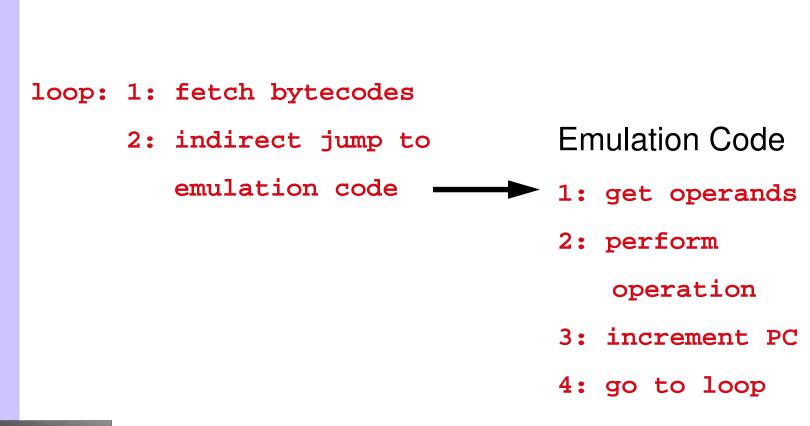
### Some instructions are complex

**lookupswitch:** "traditional" switch statement

byte 1	byte 2	byte 3	byte 4		
opcode (171)	03 byte padding				
default offset					
numbers of pairs that follow (N)					
match 1					
jump offset 1					
match 2					
jump offset 2					
•••					
match N					
jump offset N					



# **Interpreter Loop**





# **JVM: Stack-Based Architecture**

# Operands typically accessed from the stack, put back on the stack

- Example integer add:
  - –Add top 2 entries in the stack and put the result on top of the stack
  - -Typical emulation on a RISC processor
    - 1: load tos
      2: load tos-1
      3: add
      4: store tos-1



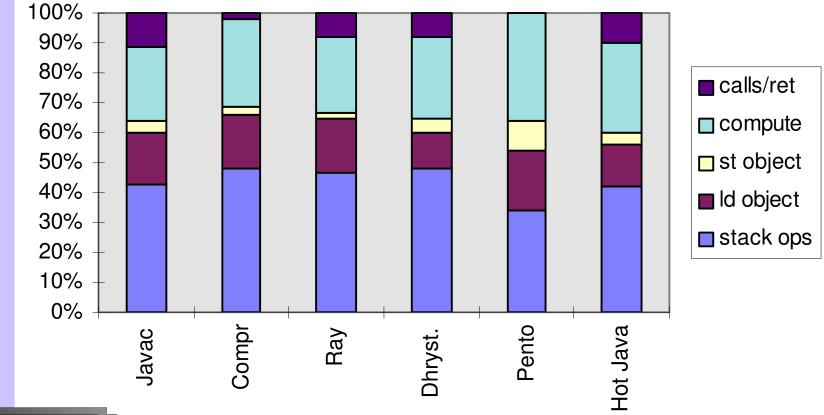
# How to Best Execute Bytecodes?

- Leverage RISC techniques developed over the past 15 years
- Implement in hardware only those instructions that make a difference

Trap for costly instructions that do not occur often
 State machines for high frequency/medium complexity instructions

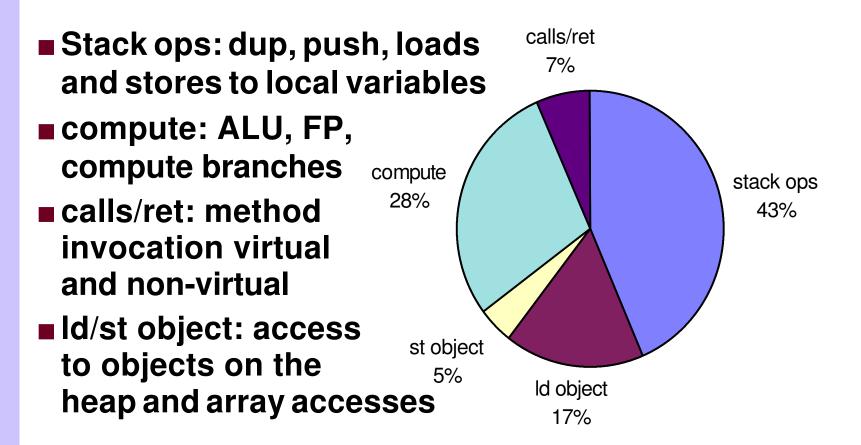


# **Dynamic Instruction Distribution**



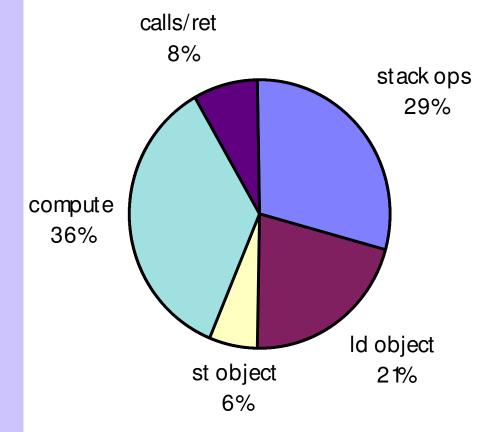


# **Composite Instruction Mix**





# **Loads from Local Variables**



- Loads from local variables move data within the chip
- Target register is often consume immediately
- Up to 60% of them can be hidden
- Resulting instruction distribution looks closer to a RISC processor



# **Pipeline Design**

### RISC pipeline attributes

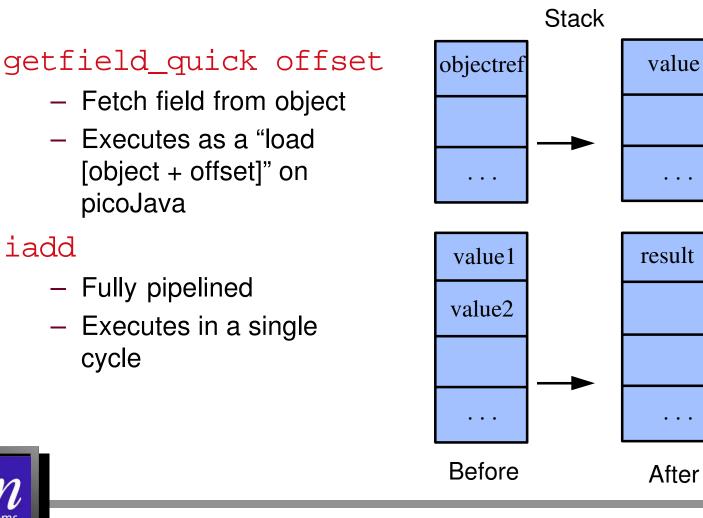
- Stages based on fundamental paths (e.g. cache access, ALU path, registers access)
- No operation on cache/memory data
- Hardwire all simple operations

### Enhance classic pipeline

- Support for method invocations
- Support for hiding loads from local variables



# Implementation of **Critical Instructions**



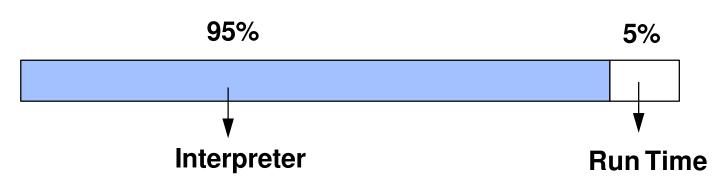


iadd

cycle

# Typical Small Benchmarks (Caffeinemarks, Pentonimo, etc.)

### Few objects, few calls, few threads

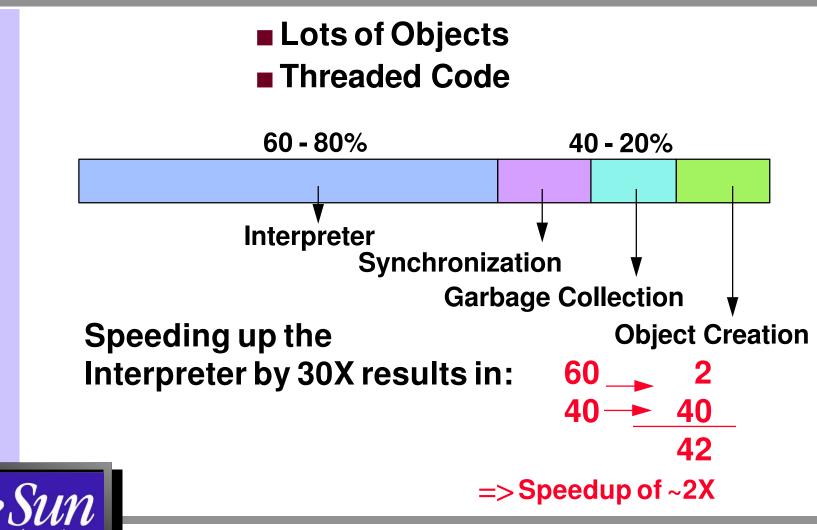


Speeding up the Interpreter by 30X results in: 95 -> 3.2 5 -> 5 8.2

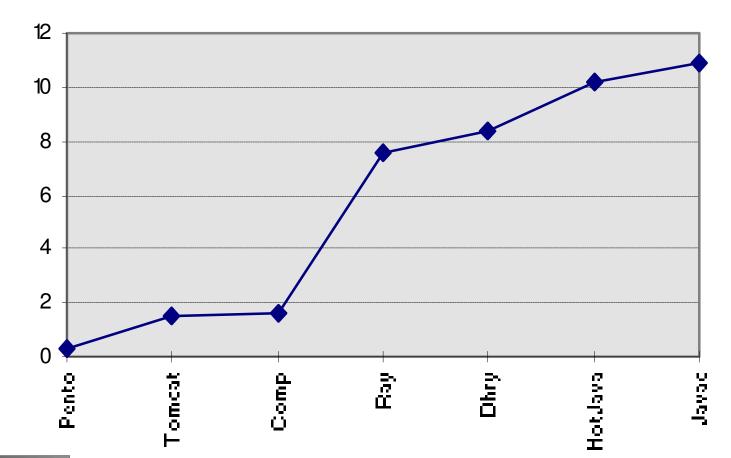
=>Speedup of ~12X



# **Representative Applications**



# **Percentage of Calls**





Varies dramatically according to benchmark type

# picoJava: A System Performance Approach

### Accelerates object-oriented programs

- simple pipeline with enhancements for features specific to bytecodes
- support for method invocation

### Accelerates runtime

(gc.c, monitor.c, threadruntime.c, etc.)

- -Support for threads
- -Support for garbage collection
- Simple but efficient, non-invasive, hardware support



# **System Programming**

- Instructions added to support system programming
  - available only "under the hood"
  - operating system functions
  - access to I/O devices
  - access to the internals of picoJava



# picoJava - Summary

Best system price/performance for running Java<sup>TM</sup>-powered applications in embedded markets

- Embedded market very sensitive to system cost and power consumption
- Interpreter and/or JIT compiler eliminated
- Excellent system performance
- Efficient implementation through use of the same methodology, process and circuit techniques developed for RISC processors

